

Logarithmic Geometry and Moduli

Tropical Geometry and Moduli Spaces

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Moduli of curves

\mathcal{M}_g - a quasiprojective variety.

Working with a non-complete moduli space is like keeping change in a pocket with holes

Angelo Vistoli

Deligne–Mumford

- $\mathcal{M}_g \subset \overline{\mathcal{M}}_g$ - moduli of *stable* curves, a modular compactification.
- allow only nodes as singularities
- What's so great about nodes?
- One answer: from the point of view of logarithmic geometry, these are the *logarithmically smooth* curves.

Logarithmic structures

Definition

A *pre logarithmic structure* is

$$X = (\underline{X}, M \xrightarrow{\alpha} \mathcal{O}_{\underline{X}}) \quad \text{or just} \quad (\underline{X}, M)$$

such that

- \underline{X} is a scheme - the *underlying scheme*
- M is a sheaf of monoids on X , and
- α is a monoid homomorphism, where the monoid structure on $\mathcal{O}_{\underline{X}}$ is the multiplicative structure.

Definition

It is a *logarithmic structure* if $\alpha : \alpha^{-1}\mathcal{O}_{\underline{X}}^* \rightarrow \mathcal{O}_{\underline{X}}^*$ is an isomorphism.

“Trivial” examples

Examples

- $(\underline{X}, \mathcal{O}_{\underline{X}}^* \hookrightarrow \mathcal{O}_{\underline{X}})$, the **trivial logarithmic** structure.
We sometimes write just \underline{X} for this structure.
- $(\underline{X}, \mathcal{O}_{\underline{X}} \xrightarrow{\sim} \mathcal{O}_{\underline{X}})$, looks as easy but surprisingly not interesting,^a and
- $(\underline{X}, \mathbb{N} \xrightarrow{\alpha} \mathcal{O}_{\underline{X}})$, where α is determined by an arbitrary choice of $\alpha(1)$.
This one is important but only pre-logarithmic.

^anot according to Dhruv!

The associated logarithmic structure

You can always fix a pre-logarithmic structure:

$$\begin{array}{ccc} \alpha^{-1}\mathcal{O}^* \hookrightarrow & M & \\ \downarrow & \vdots & \searrow \alpha \\ \mathcal{O}^* \dashrightarrow & M^a & \\ & \searrow \alpha^a & \\ & & \mathcal{O} \end{array}$$

Key examples

Example (Divisorial logarithmic structure)

Let $\underline{X}, D \subset \underline{X}$ be a variety with a divisor. We define $M_D \hookrightarrow \mathcal{O}_{\underline{X}}$:

$$M_D(U) = \left\{ f \in \mathcal{O}_{\underline{X}}(U) \mid f_{U \setminus D} \in \mathcal{O}_{\underline{X}}^\times(U \setminus D) \right\}.$$

This is particularly important for normal crossings divisors and toric divisors - these will be logarithmically smooth structures.

Example (Standard logarithmic point)

Let k be a field,

$$\begin{aligned}\mathbb{N} \oplus k^\times &\rightarrow k \\ (n, z) &\mapsto z \cdot 0^n\end{aligned}$$

defined by sending $0 \mapsto 1$ and $n \mapsto 0$ otherwise.

Works with P a monoid with $P^\times = 0$, giving the *P -logarithmic point*. This is what you get when you restrict the structure on an affine toric variety associated to P to the maximal ideal generated by $\{p \neq 0\}$.

Morphisms

A morphism of (pre)-logarithmic schemes $f : X \rightarrow Y$ consists of

- $\underline{f} : \underline{X} \rightarrow \underline{Y}$
- A homomorphism f^\flat making the following diagram commutative:

$$\begin{array}{ccc} M_X & \xleftarrow{f^\flat} & \underline{f}^{-1} M_Y \\ \alpha_X \downarrow & & \downarrow \alpha_Y \\ \mathcal{O}_X & \xleftarrow{f^\sharp} & \underline{f}^{-1} \mathcal{O}_Y \end{array}$$

Definition (Inverse image)

Given $\underline{f} : \underline{X} \rightarrow \underline{Y}$ and $Y = (\underline{Y}, M_Y)$ define the *pre-logarithmic inverse image* by composing

$$\underline{f}^{-1} M_Y \rightarrow \underline{f}^{-1} \mathcal{O}_Y \xrightarrow{f^\#} \mathcal{O}_X$$

and then the *logarithmic inverse image* is defined as

$$\underline{f}^* M_Y = (\underline{f}^{-1} M_Y)^a.$$

This is the universal logarithmic structure on \underline{X} with commutative

$$\begin{array}{ccc} (\underline{X}, \underline{f}^* M_Y) & \longrightarrow & Y \\ \downarrow & & \downarrow \\ \underline{X} & \longrightarrow & \underline{Y} \end{array}$$

$X \rightarrow Y$ is **strict** if $M_X = \underline{f}^* M_Y$.

Definition (Fibered products)

The fibered product $X \times_Z Y$ is defined as follows:

- $\underline{X \times_Z Y} = \underline{X} \times_{\underline{Z}} \underline{Y}$
- If N is the pushout of

$$\begin{array}{ccc} & \pi_Z^{-1} M_Z & \\ & \swarrow & \searrow \\ \pi_X^{-1} M_X & & \pi_Y^{-1} M_Y \end{array}$$

then the log structure on $X \times_Z Y$ is defined by N^a .

Definition (The spectrum of a Monoid algebras)

Let P be a monoid, R a ring. We obtain a monoid algebra $R[P]$ and a scheme $\underline{X} = \text{Spec } R[P]$. There is an evident monoid homomorphism $P \rightarrow R[P]$ inducing sheaf homomorphism $P_X \rightarrow \mathcal{O}_{\underline{X}}$, a pre-logarithmic structure, giving rise to a logarithmic structure

$$(P_X)^a \rightarrow \mathcal{O}_{\underline{X}}.$$

This is a basic example. It deserves a notation:

$$X = \text{Spec}(P \rightarrow R[P]).$$

The most basic example is $X_0 = \text{Spec}(P \rightarrow \mathbb{Z}[P])$.

The morphism $\underline{f} : \text{Spec}(R[P]) \rightarrow \text{Spec}(\mathbb{Z}[P])$ gives

$$X = \underline{X} \times_{\underline{X}_0} X_0.$$

Charts

A *chart* for X is given by a monoid P and a sheaf homomorphism $P_X \rightarrow \mathcal{O}_X$ to which X is associated.

This is the same as a strict morphism $X \rightarrow \text{Spec}(P \rightarrow \mathbb{Z}[P])$

Given a morphism of logarithmic schemes $f : X \rightarrow Y$, a *chart for f* is a triple

$$(P_X \rightarrow M_X, Q_Y \rightarrow M_Y, Q \rightarrow P)$$

such that

- $P_X \rightarrow M_X$ and $Q_Y \rightarrow M_Y$ are charts for M_X and M_Y , and
- the diagram

$$\begin{array}{ccc} Q_X & \longrightarrow & f^{-1}M_Y \\ \downarrow & & \downarrow \\ P_X & \longrightarrow & M_X \end{array}$$

is commutative.

Types of logarithmic structures

- We say that (\underline{X}, M_X) is *coherent* if *étale locally* at every point there is a *finitely generated* monoid P and a local chart $P_X \rightarrow \mathcal{O}_X$ for X .
- A monoid P is *integral* if $P \rightarrow P^{\text{gp}}$ is injective.
- It is *saturated* if integral and whenever $p \in P^{\text{gp}}$ and $m \cdot p \in P$ for some integer $m > 0$ then $p \in P$. I.e., not like $\{0, 2, 3, \dots\}$.
- We say that a logarithmic structure is *fine* if it is *coherent* with local charts $P_X \rightarrow \mathcal{O}_X$ with P *integral*.
- We say that a logarithmic structure is *fine and saturated* (or *fs*) if it is coherent with local charts $P_X \rightarrow \mathcal{O}_X$ with P *integral and saturated*.

Definition (The characteristic sheaf)

Given a logarithmic structure $X = (\underline{X}, M)$, the quotient sheaf $\overline{M} := M/\mathcal{O}_{\underline{X}}^\times$ is called the *characteristic sheaf* of X .

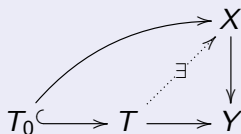
The characteristic sheaf records the combinatorics of a logarithmic structure, especially for fs logarithmic structures.

Smoothness

Definition

We define a morphism $X \rightarrow Y$ of *fine logarithmic schemes* to be *logarithmically smooth* if

- 1 $\underline{X} \rightarrow \underline{Y}$ is locally of finite presentation, and
- 2 For T_0 fine and affine and $T_0 \subset T$ strict square-0 embedding, given



there exists a lifting as indicated.

The morphism is *logarithmically étale* if the lifting in (2) is unique.

Strict smooth morphisms

Lemma

If $X \rightarrow Y$ is *strict* and $\underline{X} \rightarrow \underline{Y}$ *smooth* then $X \rightarrow Y$ is *logarithmically smooth*.

Proof.

There is a lifting

$$\begin{array}{ccccc} & & & & \underline{X} \\ & & & \nearrow & \downarrow \\ & & & \exists & \underline{Y} \\ & & & \nearrow & \\ \underline{T}_0 & \hookrightarrow & \underline{T} & \longrightarrow & \underline{Y} \end{array}$$

since $\underline{X} \rightarrow \underline{Y}$ smooth, and the lifting of morphism of monoids comes by the universal property of pullback.



Combinatorially smooth morphisms

Proposition

Say P, Q are finitely generated **integral** monoids, R a ring, $Q \rightarrow P$ a monoid homomorphism.

Write $X = \text{Spec}(P \rightarrow R[P])$ and $Y = \text{Spec}(Q \rightarrow R[Q])$.

Assume

- $\text{Ker}(Q^{\text{gp}} \rightarrow P^{\text{gp}})$ is finite and with order invertible in R ,
- $\text{TorCoker}(Q^{\text{gp}} \rightarrow P^{\text{gp}})$ has order invertible in R .

Then $X \rightarrow Y$ is logarithmically smooth.

If also the cokernel is finite then $X \rightarrow Y$ is logarithmically étale.

Key examples

- Dominant toric morphisms
- Nodal curves.
- Marked nodal curves.
- $\text{Spec } \mathbb{C}[t] \rightarrow \text{Spec } \mathbb{C}[s]$ given by $s = t^2$
- $\text{Spec } \mathbb{C}[x, y] \rightarrow \text{Spec } \mathbb{C}[t]$ given by $t = x^m y^n$
- $\text{Spec } \mathbb{C}[x, y] \rightarrow \text{Spec } \mathbb{C}[x, z]$ given by $z = xy$.
- $\text{Spec}(\mathbb{N} \rightarrow \mathbb{C}[\mathbb{N}]) \rightarrow \text{Spec}((\mathbb{N} \setminus 1) \rightarrow \mathbb{C}[(\mathbb{N} \setminus 1)])$.

Integral morphisms

Disturbing feature: the last two examples are not flat. Which ones are flat?

We define a monoid homomorphism $Q \rightarrow P$ to be *integral* if

$$\mathbb{Z}[Q] \rightarrow \mathbb{Z}[P]$$

is flat.¹

A morphism $f : X \rightarrow Y$ of logarithmic schemes is *integral* if for every geometric point x of X the homomorphism

$$(f^{-1}\overline{M}_Y)_x \rightarrow (\overline{M}_X)_x$$

of characteristic sheaves is integral.

¹Has natural universal property

Characterization of logarithmic smoothness

Theorem (K. Kato)

X, Y *fine*, $Q_Y \rightarrow M_Y$ *a chart*.

Then $X \rightarrow Y$ is *logarithmically smooth* iff there are extensions to local charts

$$(P_X \rightarrow M_X, Q_Y \rightarrow M_Y, Q \rightarrow P)$$

for $X \rightarrow Y$ such that

- $Q \rightarrow P$ combinatorially smooth, and
- $\underline{X} \rightarrow \underline{Y} \times_{\text{Spec } \mathbb{Z}[Q]} \text{Spec } \mathbb{Z}[P]$ is smooth

One direction:

$$\begin{array}{ccccc} \underline{X} & \longrightarrow & \underline{Y} \times_{\text{Spec } \mathbb{Z}[Q]} \text{Spec } \mathbb{Z}[P] & \longrightarrow & \text{Spec } \mathbb{Z}[P] \\ & & \downarrow & & \downarrow \\ & & \underline{Y} & \longrightarrow & \text{Spec } \mathbb{Z}[Q] \end{array}$$

Deformations

Proposition (K. Kato)

If $X_0 \rightarrow Y_0$ is logarithmically smooth, $Y_0 \subset Y$ a strict square-0 extension, then locally X_0 can be lifted to a smooth $X \rightarrow Y$.

Sketch of proof: locally $X_0 \rightarrow X'_0 \rightarrow Y_0$, where

$$X'_0 = Y_0 \times_{\mathrm{Spec} \mathbb{Z}[Q]} \mathrm{Spec} \mathbb{Z}[P].$$

So $X'_0 \rightarrow Y_0$ is combinatorially smooth, and automatically provided a deformation to

$$Y \times_{\mathrm{Spec} \mathbb{Z}[Q]} \mathrm{Spec} \mathbb{Z}[P],$$

and $X_0 \rightarrow X'_0$ is strict and smooth so deforms by the classical result.

Kodaira-Spencer theory

Theorem (K. Kato)

Let Y_0 be artinian, $Y_0 \subset Y$ a strict square-0 extension with ideal J , and $f_0 : X_0 \rightarrow Y_0$ *logarithmically smooth*. Then

- There is a canonical element $\omega \in H^2(X_0, T_{X_0/Y_0} \otimes f_0^* J)$ such that a logarithmically smooth deformation $X \rightarrow Y$ exists if and only if $\omega = 0$.
- If $\omega = 0$, then isomorphism classes of such $X \rightarrow Y$ correspond to elements of a torsor under $H^1(X_0, T_{X_0/Y_0} \otimes f_0^* J)$.
- Given such deformation $X \rightarrow Y$, its automorphism group is $H^0(X_0, T_{X_0/Y_0} \otimes f_0^* J)$.

Corollary

Logarithmically smooth curves are unobstructed.

Saturated morphisms

Recall that the monoid homomorphism $\mathbb{N} \xrightarrow{\cdot 2} \mathbb{N}$ gives an integral logarithmically étale map with non-reduced fibers.

Definition

- An integral $Q \rightarrow P$ of saturated monoids is said to be *saturated* if $\text{Spec}(P \rightarrow \mathbb{Z}[P]) \rightarrow \text{Spec}(Q \rightarrow \mathbb{Z}[Q])$ has reduced fibers.^a
- An integral morphism $X \rightarrow Y$ of fs logarithmic schemes is *saturated* if it has a saturated chart.

^aHas natural universal property

This guarantees that if $X \rightarrow Y$ is logarithmically smooth, then the fibers are reduced.

Log curves

Definition

A *log curve* is a morphism $f : X \rightarrow S$ of fs logarithmic schemes satisfying:

- f is logarithmically smooth,
- f is integral, i.e. flat,
- f is saturated, i.e. has reduced fibers, and
- the fibers are curves i.e. pure dimension 1 schemes.

Theorem (F. Kato)

Assume $\pi : X \rightarrow S$ is a log curve. Then

- *Fibers have at most nodes as singularities*
- *étale locally on S we can choose disjoint sections $s_i : \underline{S} \rightarrow \underline{X}$ in the nonsingular locus \underline{X}_0 of $\underline{X}/\underline{S}$ such that*
 - ▶ *Away from s_i we have that $X^0 = \underline{X}_0 \times_{\underline{S}} S$, so π is strict away from s_i*
 - ▶ *Near each s_i we have a strict étale*

$$X^0 \rightarrow S \times \mathbb{A}^1$$

with the standard divisorial logarithmic structure on \mathbb{A}^1 .

- ▶ *étale locally at a node $xy = f$ the log curve X is the pullback of*

$$\mathrm{Spec}(\mathbb{N}^2 \rightarrow \mathbb{Z}[\mathbb{N}^2]) \rightarrow \mathrm{Spec}(\mathbb{N} \rightarrow \mathbb{Z}[\mathbb{N}])$$

where $\mathbb{N} \rightarrow \mathbb{N}^2$ is the diagonal. Here the image of $1 \in \mathbb{N}$ in \mathcal{O}_S is f and the generators of \mathbb{N}^2 map to x and y .

Stable log curves

Definition

A *stable log curve* $X \rightarrow S$ is:

- a log curve $X \rightarrow S$,
- sections $s_i : \underline{S} \rightarrow \underline{X}$ for $i = 1, \dots, n$,

such that

- $(\underline{X} \rightarrow \underline{S}, s_i)$ is stable,
- the log structure is strict away from sections and singularities of fibers, and “divisorial along the sections”.

Moduli of stable log curves

We define a category $\overline{\mathcal{M}}_{g,n}^{\log}$ of stable log curves: objects are log (g, n) -curves $X \rightarrow S$ and arrows are fiber diagrams compatible with sections

$$\begin{array}{ccc} X_1 & \longrightarrow & X_2 \\ \downarrow & & \downarrow \\ S_2 & \longrightarrow & S_2 \end{array}$$

There is a forgetful functor

$$\begin{array}{ccc} \overline{\mathcal{M}}_{g,n}^{\log} & \longrightarrow & \mathcal{L}\text{og}\mathcal{S}\text{ch}^{\text{fs}} \\ (X \rightarrow S) & \mapsto & S. \end{array}$$

So $\overline{\mathcal{M}}_{g,n}^{\log}$ is a category fibered in groupoids over $\mathcal{L}\text{og}\mathcal{S}\text{ch}^{\text{fs}}$.

Moduli of stable log curves (continued)

We also have a forgetful functor

$$\begin{array}{ccc} \overline{\mathcal{M}}_{g,n}^{\log} & \longrightarrow & \overline{\mathcal{M}}_{g,n} \\ (X \rightarrow S) & \mapsto & (\underline{X} \rightarrow \underline{S}) \end{array}$$

Note that the Deligne–Knudsen–Mumford moduli stack $\overline{\mathcal{M}}_{g,n}$ has a natural logarithmic smooth structure $M_{\Delta_{g,n}}$ given by the boundary divisor. As such it represents a category fibered in groupoids $(\overline{\mathcal{M}}_{g,n}, M_{\Delta_{g,n}})$ over $\mathcal{L}\text{ogS}\text{ch}^{\text{fs}}$.

Theorem (F. Kato)

$$\overline{\mathcal{M}}_{g,n}^{\log} \simeq (\overline{\mathcal{M}}_{g,n}, M_{\Delta_{g,n}}).$$

The stack $\overline{\mathcal{M}}_{g,n}^{\log}$ is huge!

- Given a stable log curve $X \rightarrow S$ and an arbitrary log point S' then $X \times S' \rightarrow S \times S'$ is a stable log curve **on the same underlying scheme**.
- Given $S = \text{Spec}(\mathbb{N}^2 \rightarrow \mathbb{C})$, a stable log curve $X \rightarrow S$ and a quotient $\mathbb{N}^2 \rightarrow \mathbb{N}$ we get a log point $S' = \text{Spec}(\mathbb{N} \rightarrow \mathbb{C}) \rightarrow S$ and a new stable log curve $X' \rightarrow S'$ **on the same underlying scheme**.
- We claim that some stable log curves $X \rightarrow S$ are more fundamental than others.
- Not all logarithmic moduli problems are so lucky! There are issues with logarithmic \mathbb{G}_m and logarithmic Picard (Molcho-Wise).

Minimality

Given a stable curve $\underline{X} \rightarrow \underline{S}$ we define

$$X^{\min} = \underline{X} \times_{\overline{\mathcal{M}}_{g,n+1}} \overline{\mathcal{M}}_{g,n+1}^{\log} \quad \text{and} \quad S^{\min} = \underline{S} \times_{\overline{\mathcal{M}}_{g,n}} \overline{\mathcal{M}}_{g,n}^{\log}.$$

The logarithmic structures $X^{\min} \rightarrow S^{\min}$ are called the *minimal* or *basic* logarithmic structures on a log curve.

We write

$$S^{\min} = (\underline{S}, M_{X/S}^S) \quad \text{and} \quad X^{\min} = (\underline{X}, M_{X/S}^X).$$

Fundamental diagram

$$\begin{array}{ccccc} X & \longrightarrow & X^{\min} & \longrightarrow & \mathcal{X}_{g,n}^{\log} \\ \downarrow & & \downarrow & & \downarrow \\ S & \longrightarrow & S^{\min} & \longrightarrow & \overline{\mathcal{M}}_{g,n}^{\log} \\ & \searrow & \downarrow & & \downarrow \\ & & \underline{S} & \longrightarrow & \overline{\mathcal{M}}_{g,n} \end{array}$$

- $\overline{\mathcal{M}}_{g,n}^{\log}$ parametrizes stable log curves over $\mathcal{L}\text{og}\mathcal{S}\text{ch}^{\text{fs}}$
- $\overline{\mathcal{M}}_{g,n}$ parametrizes **minimal** stable log curves over $\mathcal{S}\text{ch}$.

Stable logarithmic maps

Definition

A **stable logarithmic map** is a diagram

$$\begin{array}{ccc} C & \xrightarrow{f} & X \\ \downarrow \pi & & \\ S & & \end{array}$$

s_i (curved arrow from S to C)

where

- $(C/S, s_i)$ is a prestable log curve, and
- in fibers $\text{Aut}(\underline{C}_s \rightarrow \underline{X}, s_i)$ is finite.

contact orders

Apart from the underlying discrete data $\underline{\Gamma} = (g, \beta, n)$, a stable logarithmic map has *contact orders* c_i at the marked points.

At each such point the logarithmic structure at C has a factor \mathbb{N} , and the contact order is the homomorphism $f^* M_X \xrightarrow{c_i} \mathbb{N}$ at that marked point.

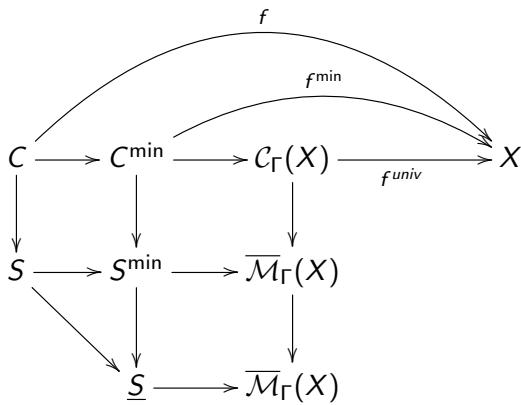
We collect the numerical data under the umbrella $\Gamma = (g, \beta, c_i)$.

Stable logarithmic maps (continued)

Theorem (Gross-Siebert, Chen, \(\infty\)-Chen-Marcus-Wise)

Let X be projective logarithmically scheme. Stable logarithmic maps to X form a logarithmic Deligne–Mumford stack $\overline{\mathcal{M}}_{\Gamma}(X)$. It is finite and representable over $\overline{\mathcal{M}}_{\Gamma}(X)$.

Fundamental diagram



We are in search of a moduli stack $\underline{\overline{\mathcal{M}}}_{\Gamma}(X)$ parametrizing *minimal* stable logarithmic maps over $\mathcal{S}ch$.

As such it comes with a logarithmic structure $\overline{\mathcal{M}}_{\Gamma}(X)$ which parametrizes *all* stable logarithmic maps over $\mathcal{L}og\mathcal{S}ch^{fs}$.

Stable logarithmic maps (continued)

This requires two steps:

- first find a morphism from $(C \rightarrow S, f : C \rightarrow X)$ to a *minimal* object $(C^{\min} \rightarrow S^{\min}, f^{\min} : C^{\min} \rightarrow X)$.
- then show that the object $(C^{\min} \rightarrow S^{\min}, f^{\min} : C^{\min} \rightarrow X)$ has a versal deformation space, $\underline{S} \rightarrow \underline{V}$, with universal family

$$(C_V^{\min} \rightarrow V^{\min}, f_V^{\min} : C_V^{\min} \rightarrow X),$$

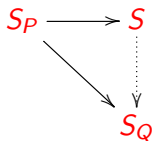
whose fibers over any $\underline{T} \rightarrow \underline{V}$ are also minimal.

Minimal stable logarithmic maps

We consider X toric and a stable logarithmic map $(C/S, f : C \rightarrow X)$ over a P -logarithmic point S .

We wish to find a minimal Q -logarithmic point and a logarithmic map over it through which our object factors.

We might as well first pull back and replace S by a standard, $P = \mathbb{N}$ -logarithmic point!



Indeed if $S_P \rightarrow S$ is such a point and $S_P \rightarrow S_Q$ is its map to the minimal structure, then minimality says there is a unique $S \rightarrow S_Q$ factoring $S_P \rightarrow S_Q$

At the generic points

The curve C has components C_i with generic points η_i corresponding to vertices in the dual graph, and nodes q_j with local equations $xy = g_j$ corresponding to edges in the dual graph.

- The map f sends η_i to some stratum X_i of X with cone σ_i having lattice $N_i \subset \sigma_i$.
- Departing from toric conventions we denote $M_i = N_i^\vee = \underline{\text{Hom}}(N_i, \mathbb{N})$.
- Since the logarithmic structure of C at η_i is the pullback of the structure on S , we have a map $f_i^\flat : M_i \rightarrow P$.
- It can dually be viewed as a map $P^\vee \rightarrow N_i$, or an element $v_i \in N_i$.
- If that were all we had, our final object would be $Q^\vee = \prod N_i$, and dually the initial monoid $Q = \oplus M_i$.

But the nodes impose crucial conditions.

At the nodes

- At a node q with branches η_q^1, η_q^2 we similarly have a map

$$f_q^b : M_q \rightarrow P \oplus^{\mathbb{N}} \mathbb{N}^2.$$

- Unfortunately it is unnatural to consider maps into a coproduct, and we give an alternate description of

$$P \oplus^{\mathbb{N}} \mathbb{N}^2 = P \langle \log x, \log y \rangle / (\log x + \log y = \rho_q)$$

where $\rho_q = \log g_q \in P$.

- Recall that the stalk of a sheaf at a point q maps, via a “generalization map”, to the stalk at any point specializing to q , such as η_q^1, η_q^2 .
- The stalk at either η_q^1, η_q^2 is P .

At the nodes (continued)

$$P \oplus^{\mathbb{N}} \mathbb{N}^2 = P \langle \log x, \log y \rangle / (\log x + \log y = \rho_q)$$

The map to the stalk P at η_q^1 where $x = 0$ sends $\log y \mapsto 0$, and so $\log x \mapsto \rho_q$.

The map to the stalk P at η_q^2 where $y = 0$ sends $\log x \mapsto 0$, and so $\log y \mapsto \rho_q$.

This means that we have a monoid homomorphism, which is clearly injective,

$$P \oplus^{\mathbb{N}} \mathbb{N}^2 \rightarrow P \times P.$$

Its image is precisely the set of pairs

$$\{(p_1, p_2) \mid p_2 - p_1 \in \mathbb{Z}\rho_q\}$$

At the nodes (continued)

$$P \oplus^{\mathbb{N}} \mathbb{N}^2 = \{(p_1, p_2) \mid p_2 - p_1 \in \mathbb{Z}\rho_q\}$$

- $f_q^b : M_q \rightarrow P \times P$,
- $(p_2 - p_1) \circ f_q^b : M_q \rightarrow \mathbb{Z}\rho_q \subset P^{\text{gp}}$.

Or better: we have $u_q : M_q \rightarrow \mathbb{Z}$ such that

$$(p_2 - p_1) \circ f_q^b(m) = u_q(m) \cdot \rho_q. \tag{0.0.1}$$

Putting nodes and generic points together

The maps $p_1 \circ f_q^b : M_q \rightarrow P$ and $p_2 \circ f_q^b : M_q \rightarrow P$, since they come from maps of sheaves, are compatible with generization maps.

$p_1 \circ f_q^b : M_q \rightarrow P$ is the composition $M_q \rightarrow M_{\eta_q^1} \rightarrow P$

$p_2 \circ f_q^b : M_q \rightarrow P$ is the composition of $M_q \rightarrow M_{\eta_q^2} \rightarrow P$

the data of $p_1 \circ f_q^b$ and $p_2 \circ f_q^b$ is already determined by the data at the generic points η_i of the curve.

Putting nodes and generic points together (continued)

The *additional* data the node provides is the element $\rho_q \in P$ and homomorphism $u_q : M_q \rightarrow \mathbb{Z}$, in such a way that equation

$$(p_2 - p_1) \circ f_q^b(m) = u_q(m) \cdot \rho_q. \quad (0.0.2)$$

holds.

$$Q_f = \left(\left(\prod_{\eta} M_{\sigma_{\eta}} \times \prod_q \mathbb{N} \right) / R \right)^{sat}$$

where the saturated submonoid R is generated by all the relations implied by equation (??)

Putting nodes and generic points together (continued)

$$Q_f = \left(\left(\prod_{\eta} M_{\sigma_{\eta}} \times \prod_q \mathbb{N} \right) / R \right)^{\text{sat}}$$

It is quite a bit more natural to describe the dual lattice

$$Q_f^{\vee} = \left\{ ((v_{\eta}), (e_q)) \in \prod_{\eta} N_{\sigma_{\eta}} \times \prod_q \mathbb{N} \mid \forall \begin{array}{ccc} \bullet & \xrightarrow{q} & \bullet \\ \eta_q^1 & & \eta_q^2 \\ v_q^1 - v_q^2 = e_q u_q & & \end{array} \right\}.$$

Tropical interpretation

Given a map f over an \mathbb{N} -point we have a graph in $\Sigma(X)$ with

- vertices $v_i \in N_{\sigma_\eta} \subset \sigma_\eta$
- edges proportional to $u_q \in N_q^{\text{gp}}$ such that $v_q^1 - v_q^2 = e_q u_q$

this means

- The equations $v_q^1 - v_q^2 = e_q u_q$ define the cone of all such graphs
- Q_f^\vee is the integer lattice in that cone.

Theorem (Gross-Siebert)

The minimal object exists, with characteristic sheaf Q_f , dual to the lattice in the corresponding space of tropical curves.