

Punctured logarithmic invariants and gluing

Dan Abramovich, Brown University

Work with **Qile Chen**, **Mark Gross** and **Bernd Siebert**

Other work by Parker, Tehrani, Dhruv, Fan-Tseng-Wu-You

Geometry and Analysis Seminar, Oxford

February 8, 2021

The gluing result

Here is a minimalist¹ statement:

Theorem (ACGS 2020)

The evaluation maps $\widetilde{\mathcal{M}}(X, \tau) \rightarrow X^n$ of the moduli stack $\widetilde{\mathcal{M}}(X, \tau)$ of *stable marked punctured curves of type τ* in a *log scheme X* are *virtually idealized log smooth*.

Given an edge of τ with splitting τ' we have a cartesian splitting diagram

$$\begin{array}{ccc} \widetilde{\mathcal{M}}(X/B, \tau) & \longrightarrow & \widetilde{\mathcal{M}}(X, \tau') \\ \downarrow & & \downarrow \\ X & \longrightarrow & X \times X \end{array}$$

of fs log stacks with compatible virtual structure.

¹cubist impressionistic

The gluing result

Here is a minimalist¹ statement:

Theorem (ACGS 2020)

The evaluation maps $\widetilde{\mathcal{M}}(X, \tau) \rightarrow X^n$ of the moduli stack $\widetilde{\mathcal{M}}(X, \tau)$ of *stable marked punctured curves of type τ* in a *log scheme X* are *virtually idealized log smooth*.

Given an edge of τ with splitting τ' we have a cartesian splitting diagram

$$\begin{array}{ccc} \widetilde{\mathcal{M}}(X/B, \tau) & \longrightarrow & \widetilde{\mathcal{M}}(X, \tau') \\ \downarrow & & \downarrow \\ X & \longrightarrow & X \times X \end{array}$$

of fs log stacks with compatible virtual structure.

There is a lot I need to explain and motivate.

¹cubist impressionistic

Rational plane curves

Definition

$$N_d = \# \left\{ \begin{array}{l} C \subset \mathbb{P}^2 \text{ a rational curve,} \\ \deg C = d, \text{ and} \\ p_1, \dots, p_{3d-1} \in C \end{array} \right\}.$$

Kontsevich's theorem

Definition

$$N_d = \# \left\{ \begin{array}{l} C \subset \mathbb{P}^2 \text{ a rational curve,} \\ \deg C = d, \text{ and} \\ p_1, \dots, p_{3d-1} \in C \end{array} \right\}.$$

Theorem (Kontsevich)

For $d > 1$ we have

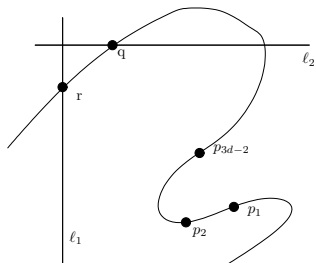
$$N_d = \sum_{\substack{d = d_1 + d_2 \\ d_1, d_2 > 0}} N_{d_1} N_{d_2} \left(d_1^2 d_2^2 \binom{3d-4}{3d_1-2} - d_1^3 d_2 \binom{3d-4}{3d_1-1} \right).$$

The first few numbers are

$$N_1 = 1, N_2 = 1, N_3 = 12, N_4 = 620, N_5 = 87304.$$

Kontsevich's theorem: setup

1-parameter family $C \rightarrow B$: fix only p_1, \dots, p_{3d-2} ,
and two lines ℓ_1, ℓ_2 meeting at a point called p_{3d-3} :

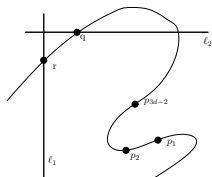


We choose q among $C \cap \ell_1$, and r among $C \cap \ell_2$.

Kontsevich's theorem: preview

The equation

$$N_d = \sum_{\substack{d = d_1 + d_2 \\ d_1, d_2 > 0}} N_{d_1} N_{d_2} \left(d_1^2 d_2^2 \binom{3d-4}{3d_1-2} - d_1^3 d_2 \binom{3d-4}{3d_1-1} \right)$$



translates to

$$\deg_B(p_1, p_2 | q, r) = \deg_B(p_1, q | p_2, r)$$

coming from the cross ratio map $\lambda : B \rightarrow \mathbb{P}^1$.

Kontsevich's theorem: computation

$$\deg_B(p_1, p_2|q, r)$$

 $=$

$$\deg_B(p_1, q|p_2, r)$$

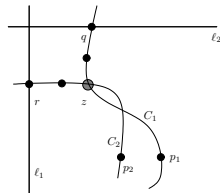
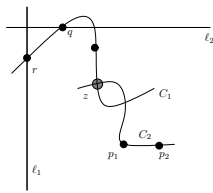
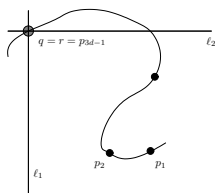


Figure: N_d

 $+$

$$d_1^3 d_2^2 \binom{3d-4}{3d_1-1} \sum N_{d_1} N_{d_2}$$

 $=$

$$\sum d_1^2 d_2^2 \binom{3d-4}{3d_1-2} N_{d_1} N_{d_2},$$

Kontsevich's theorem: computation

$$\deg_B(p_1, p_2|q, r)$$

=

$$\deg_B(p_1, q|p_2, r)$$

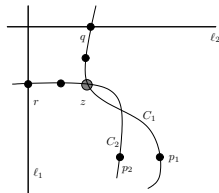
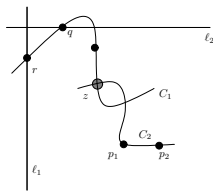
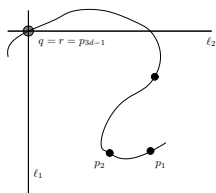


Figure: N_d

+

$$d_1^3 d_2^2 \binom{3d-4}{3d_1-1} \sum N_{d_1} N_{d_2}$$

=

$$\sum d_1^2 d_2^2 \binom{3d-4}{3d_1-2} N_{d_1} N_{d_2},$$

as needed!



Moduli spaces

- Let $\overline{\mathcal{M}}(X, \tau)$ be the Kontsevich moduli stack of stable maps in X
- with type specified by decorated graph $\tau = (G, h, \beta)$.
- to each vertex v of G we assign a genus $h(v)$ and a curve class $\beta(v)$.
- The legs are marked:

The gluing principle on moduli spaces

Proposition

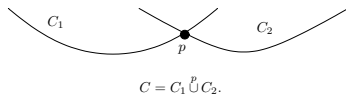
The evaluation maps $\overline{\mathcal{M}}(X, \tau) \rightarrow X^m$ are *virtually smooth*.

Given an edge of τ with splitting τ' we have a cartesian splitting diagram

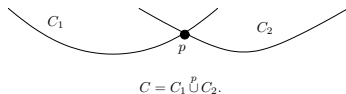
$$\begin{array}{ccc} \overline{\mathcal{M}}(X/B, \tau) & \longrightarrow & \overline{\mathcal{M}}(X, \tau') \\ \downarrow & & \downarrow \\ X & \longrightarrow & X \times X \end{array}$$

of stacks with compatible virtual fundamental classes.

The gluing principle on curves



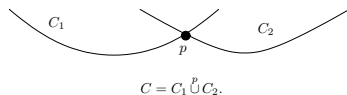
The gluing principle on curves



implying

$$\begin{aligned} \text{Hom}(C, X) &= \text{Hom}(C_1, X) \times_{\text{Hom}(p, X)} \text{Hom}(C_2, X) \\ &= \text{Hom}(C_1, X) \times_X \text{Hom}(C_2, X) \end{aligned}$$

The gluing principle on curves



implying

$$\begin{aligned} \text{Hom}(C, X) &= \text{Hom}(C_1, X) \times_{\text{Hom}(p, X)} \text{Hom}(C_2, X) \\ &= \text{Hom}(C_1, X) \times_X \text{Hom}(C_2, X) \end{aligned}$$

spreading out to

$$\begin{array}{ccc} \overline{\mathcal{M}}(X/B, \tau) & \longrightarrow & \overline{\mathcal{M}}(X, \tau') \\ \downarrow & & \downarrow \\ X & \longrightarrow & X \times X \end{array}$$

From gluing to quantum cohomology

Note that we relied on

$$X = \text{Hom}(p, X).$$

One defines quantum cohomology based on the operation

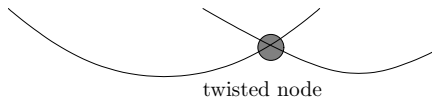
$$\langle \gamma_1 \dots \gamma_n \rangle_\beta = e_{n+1*} ([M]^{\text{virt}} \cap e_1^* \gamma_1 \dots e_n^* \gamma_n).$$

Associativity is a result of gluing.

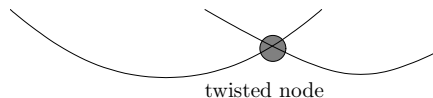
Enter orbifolds

- Let \mathcal{X} be a smooth proper DM stack (an algebraic orbifold).
- Let $\overline{\mathcal{M}}(\mathcal{X}, \tau)$ be the moduli stack of stable maps in \mathcal{X}
- with type specified by decorated graph τ .
- to each vertex v of G we assign a genus $h(v)$ and a curve class $\beta(v)$.
- **The legs are marked by inertia components.**
- We insist that every marking is a section of \mathcal{C} .

The gluing principle on orbifold curves



The gluing principle on orbifold curves



implying

$$\begin{aligned} \text{Hom}(\mathcal{C}, \mathcal{X}) &= \text{Hom}(\mathcal{C}_1, \mathcal{X}) \times_{\text{Hom}(\Sigma, \mathcal{X})} \text{Hom}(\mathcal{C}_2, \mathcal{X}) \\ &= \text{Hom}(\mathcal{C}_1, \mathcal{X}) \times_{\mathcal{IX}} \text{Hom}(\mathcal{C}_2, \mathcal{X}) \end{aligned}$$

The map on the right is the evaluation map e_2 . The map on the left is the **twisted evaluation map \check{e}_1** .

Orbifold gluing

Proposition

The evaluation maps $\overline{\mathcal{M}}(\mathcal{X}, \tau) \rightarrow \mathcal{I}\mathcal{X}^m$ are *virtually smooth*.

Given an edge of τ with splitting τ' we have a cartesian splitting diagram

$$\begin{array}{ccc} \overline{\mathcal{M}}(\mathcal{X}, \tau) & \longrightarrow & \overline{\mathcal{M}}(\mathcal{X}, \tau') \\ \downarrow & & \downarrow \\ \mathcal{I}\mathcal{X} & \longrightarrow & \mathcal{I}\mathcal{X} \times \mathcal{I}\mathcal{X} \end{array}$$

of stacks with compatible virtual fundamental classes.

From orbifold gluing to orbifold quantum cohomology

Note that we relied on

$$\mathcal{IX} = \text{Hom}(\Sigma, \mathcal{X}).$$

One defines quantum cohomology on $H^*(\mathcal{IX})$ based on the operation

$$\langle \gamma_1 \cdots \gamma_n \rangle_{\beta}^{\mathcal{X}} = \check{e}_{n+1*} ([M]^{\text{virt}} \cap e_1^* \gamma_1 \cdots e_n^* \gamma_n).$$

Associativity is a result of gluing.

Log geometry (K. Kato, Fontaine, Illusie; Ogus)

- **schemes** are glued from closed subsets of **affine spaces** - the standard-issue smooth spaces.

Log geometry (K. Kato, Fontaine, Illusie; Ogus)

- **schemes** are glued from closed subsets of **affine spaces** - the standard-issue smooth spaces.
- **log schemes** are étale glued from closed subsets of **affine toric varieties** - the standard-issue log smooth spaces.

Log geometry (K. Kato, Fontaine, Illusie; Ogus)

- **schemes** are glued from closed subsets of **affine spaces** - the standard-issue smooth spaces.
- **log schemes** are étale glued from closed subsets of **affine toric varieties** - the standard-issue log smooth spaces.
- **idealized log schemes** are étale glued from closed subsets of **monomial subschemes of affine toric varieties** - the standard-issue idealized log smooth spaces.

Log structures (K. Kato, Fontaine–Illusie)

- a **log structure** is a monoid homomorphism $\alpha : M \rightarrow \mathcal{O}_X$
- such that $\alpha^* \mathcal{O}^\times \rightarrow \mathcal{O}^\times$ is an isomorphism.

Log structures (K. Kato, Fontaine–Illusie)

- a **log structure** is a monoid homomorphism $\alpha : M \rightarrow \mathcal{O}_X$
- such that $\alpha^* \mathcal{O}^\times \rightarrow \mathcal{O}^\times$ is an isomorphism.
- Morphisms are given by natural commutative diagrams. . .

Log structures (K. Kato, Fontaine–Illusie)

- a **log structure** is a monoid homomorphism $\alpha : M \rightarrow \mathcal{O}_X$
- such that $\alpha^* \mathcal{O}^\times \rightarrow \mathcal{O}^\times$ is an isomorphism.
- Morphisms are given by natural commutative diagrams. . .
- A key example is the log structure associated to an open $U \subset X$,
- where $M = \mathcal{O}_X \cap \mathcal{O}_U^\times$.

Idealized log structures (Ogus)

- a **idealized log structure** is a log structure $\alpha : M \rightarrow \mathcal{O}_X$
- along with a monoid ideal $K \subset M$,
- such that $\alpha(K) = 0 \in \mathcal{O}$.

Toric and log smooth Log structures (K. Kato)

- When X is a toric variety and U the torus this is a prototypical example of a log smooth structure.
- In this case the monoid is associated to the regular monomials, with \mathcal{O}^\times thrown in.

Toric and log smooth Log structures (K. Kato)

- When X is a toric variety and U the torus this is a prototypical example of a log smooth structure.
- In this case the monoid is associated to the regular monomials, with \mathcal{O}^\times thrown in.
- In general X is log smooth if it is étale locally toric.
- A morphism $X \rightarrow Y$ is log smooth if it is étale locally a base change of a dominant morphism of toric varieties.

Log curves

- A **log curve** is a reduced 1-dimensional fiber of a flat **log smooth morphism**.
- F. Kato showed that these are the same as nodal marked curves, with “the natural” log structure.
- A **punctured curve** is the idealized version of the above.

Log curves under the microscope

- Say $C \rightarrow S$ a log curve, $S = \text{Spec}(M_S \rightarrow k)$.

Log curves under the microscope

- Say $C \rightarrow S$ a log curve, $S = \text{Spec}(M_S \rightarrow k)$.
- A general point of C looks like $\text{Spec}(M_S \rightarrow k[x])$.

Log curves under the microscope

- Say $C \rightarrow S$ a log curve, $S = \text{Spec}(M_S \rightarrow k)$.
- A general point of C looks like $\text{Spec}(M_S \rightarrow k[x])$.
- A node looks like $\text{Spec}(M \rightarrow k[x, y]/(xy))$, where

$$M = M_S \langle \log x, \log y \rangle / (\log x + \log y = \log t), t \in M_S.$$

Log curves under the microscope

- Say $C \rightarrow S$ a log curve, $S = \text{Spec}(M_S \rightarrow k)$.
- A general point of C looks like $\text{Spec}(M_S \rightarrow k[x])$.
- A node looks like $\text{Spec}(M \rightarrow k[x, y]/(xy))$, where

$$M = M_S \langle \log x, \log y \rangle / (\log x + \log y = \log t), t \in M_S.$$

- A marked point looks like $\text{Spec}(M \rightarrow k[x])$ where

$$M = M_S \oplus \mathbb{N} \log x.$$

Punctured curves under the microscope

- A **puncturing** of a marked curve is a log structure M at a marked point with

$$M_S + \mathbb{N} \log x \subseteq M \subsetneq M_S + \mathbb{Z} \log x.$$

- It is an instance of an **idealized log smooth** scheme.
- In particular the **splitting of a node is a punctured curve**.
- In what follow, I insist that every marking is given with a section.

Splitting

- Consider $X \rightarrow \mathbb{A}^1$ the total space of $xy = t$, and
- $C \rightarrow S$ given by $\{y = 0\} \rightarrow \{t = 0\}$.
- At the origin $M_S + \mathbb{N} \log x \subsetneq M \subsetneq M_S + \mathbb{Z} \log x$.
- It is not a log curve, but rather a **punctured curve**.

Stable punctured log maps

- Fix X a nice log smooth scheme. It has a cone complex $\Sigma(X)$ with integer lattice.
- A **stable punctured log map** $C \rightarrow X$ is a log morphism with stable underlying morphism of schemes.

Stable punctured log maps

- Fix X a nice log smooth scheme. It has a cone complex $\Sigma(X)$ with integer lattice.
- A **stable punctured log map** $C \rightarrow X$ is a log morphism with stable underlying morphism of schemes.
- Marked points record **contact orders** with divisors of X .
- These are recorded by integer vectors living in the space $\Sigma(X)(\mathbb{N})$.

Stable punctured log maps

- Fix X a nice log smooth scheme. It has a cone complex $\Sigma(X)$ with integer lattice.
- A **stable punctured log map** $C \rightarrow X$ is a log morphism with stable underlying morphism of schemes.
- Marked points record **contact orders** with divisors of X .
- These are recorded by integer vectors living in the space $\Sigma(X)(\mathbb{N})$.
- Stable punctured log maps have “standard issue” log structure, called **minimal**.

Stable punctured log maps

- Fix X a nice log smooth scheme. It has a cone complex $\Sigma(X)$ with integer lattice.
- A **stable punctured log map** $C \rightarrow X$ is a log morphism with stable underlying morphism of schemes.
- Marked points record **contact orders** with divisors of X .
- These are recorded by integer vectors living in the space $\Sigma(X)(\mathbb{N})$.
- Stable punctured log maps have “standard issue” log structure, called **minimal**.

Theorem ([ACGS])

$\mathcal{M}(X, \tau)$, the stack of minimal stable punctured log maps of **type** τ , is a Deligne–Mumford stack which is finite and representable over $\mathcal{M}(\underline{X}, \underline{\tau})$.

Prestable and cut maps (B. Parker)

- There is a range of choices for the punctured structure.

Prestable and cut maps (B. Parker)

- There is a range of choices for the punctured structure.
- For moduli of maps purposes, we use **prestable structures**:
- It is the minimal puncturing accommodating the map.
- For the purpose of gluing along sections, one can use **cut curve structures**.
- There is the maximal puncturing accommodating the section.
- The resulting categories are equivalent.

Tropical picture

- X has a cone complex $\Sigma(X)$ with integer lattice.
- $C \rightarrow S$ has cone complex $\Sigma(C) \rightarrow \Sigma(S)$. The fiber over $u \in \Sigma(S)$ is a **tropical curve**:
- Components give vertices, nodes give edges, and punctured points give legs.
- Usual marked points give infinite legs.
- Truly punctured points (and cut curves) give finite legs.
- A stable punctured log map gives $\Sigma(C) \rightarrow \Sigma(X)$, a family of tropical curves in $\Sigma(X)$.
- The sections mark the legs.
- Minimality is beautifully encoded in this picture. . .

An Analogy: Orbifold vs. Logarithmic cohomology

- If \mathcal{X} is an orbifold, **Chen-Ruan** defined orbifold and quantum cohomology based on $H^*(\bar{\mathcal{I}}(\mathcal{X}), \mathbb{Q})$.
- $\bar{\mathcal{I}}(\mathcal{X})$, the **rigidified inertia stack** is the moduli space of orbifold points in \mathcal{X} ,
- whose components, **twisted sectors**, correspond to (x, ϕ) where $x \in \mathcal{X}$ and $\phi \in \text{Aut}(x)$.

An Analogy: Orbifold vs. Logarithmic cohomology

- If \mathcal{X} is an orbifold, **Chen-Ruan** defined orbifold and quantum cohomology based on $H^*(\bar{\mathcal{I}}(\mathcal{X}), \mathbb{Q})$.
- $\bar{\mathcal{I}}(\mathcal{X})$, the **rigidified inertia stack** is the moduli space of orbifold points in \mathcal{X} ,
- whose components, **twisted sectors**, correspond to (x, ϕ) where $x \in \mathcal{X}$ and $\phi \in \text{Aut}(x)$.
- Chen Ruan cohomology pairs ϕ with ϕ^{-1} .

An Analogy: Orbifold vs. Logarithmic cohomology

- If \mathcal{X} is an orbifold, **Chen-Ruan** defined orbifold and quantum cohomology based on $H^*(\bar{\mathcal{I}}(\mathcal{X}), \mathbb{Q})$.
- $\bar{\mathcal{I}}(\mathcal{X})$, the **rigidified inertia stack** is the moduli space of orbifold points in \mathcal{X} ,
- whose components, **twisted sectors**, correspond to (x, ϕ) where $x \in \mathcal{X}$ and $\phi \in \text{Aut}(x)$.
- Chen Ruan cohomology pairs ϕ with ϕ^{-1} .
- If X is a log scheme, **Gross-Hacking-Keel-Siebert...** define the ring of theta functions,
- based on the moduli space $\mathcal{P}(X)$ of **log points** in \mathcal{X} ,
- whose components correspond to (x, u) where $x \in X$ and u a contact order at x , namely $u \in \Sigma(X)(\mathbb{N})$.
- **Gluing pairs u and $-u$.**

Gluing punctured curves

Lemma

Let C_1°, C_2° be two *cut curves* with underlying curves \underline{C}_i , over a log scheme W with sections $W \rightarrow C_i^\circ$ along the puncture.

There is a unique log structure C , log smooth over W on the nodal curve $\underline{C} = \underline{C}_1 \cup^P \underline{C}_2$, with a section at the node, restricting to C_i° . Moreover, C has the coproduct property:

$$\mathrm{Hom}(C, X) = \mathrm{Hom}(C_1^\circ, X) \times_{\mathrm{Hom}(W, X)} \mathrm{Hom}(C_2^\circ, X).$$

This has a slightly more involved implication in terms of pre-stable punctured maps.

Gluing punctured curves: moduli

This gives part of the first claim:

Theorem (ACGS 2020)

... *The following is cartesian:*

$$\begin{array}{ccc} \widetilde{\mathcal{M}}(X/B, \tau) & \longrightarrow & \widetilde{\mathcal{M}}(X, \tau') \\ \downarrow & & \downarrow \\ X & \longrightarrow & X \times X \end{array}$$

The end

Thank you for your attention