

Semistable reduction - a progress report

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Moduli and Hodge Theory
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Outline

- Statement of two results [NTW], [ALT],
- The one-dimensional base case [KKMS]
- Old results and conjectures over larger base [NK]
- Relative desingularization in the age of log stacks

This lays groundwork for Tëmkin's lecture tomorrow.

Relatively functorial toroidalization (\aleph -Tëmkin-Włodarczyk)

Theorem (\aleph TW p2020)

Let $X \rightarrow B$ be a dominant morphism of complex varieties. There is a *relatively functorial* diagram

$$\begin{array}{ccc} X' & \rightarrow & X \\ \downarrow & & \downarrow \\ B' & \rightarrow & B \end{array}$$

with

- $B' \rightarrow B$ and $X' \rightarrow B'$ modifications,
- $X' \rightarrow B'$ *logarithmically smooth*.
- In particular,
 - ▶ if the generic fiber of $X \rightarrow B$ is smooth it is not modified, and
 - ▶ a group actions along the fibers of $X \rightarrow B$ lifts to X' .

Semistable reduction (Adiprasito-Liu-Tëmkin)

Theorem (ALT p2018)

Let $X \rightarrow B$ be a generically smooth complex projective family of varieties.
There is a diagram

$$\begin{array}{ccc} X_1 & \rightarrow & X \\ \downarrow & & \downarrow \\ B_1 & \rightarrow & B \end{array}$$

with

- $B_1 \rightarrow B$ an alteration,
- $X_1 \rightarrow (X \times_B B_1)_{\text{main}}$ a modification of the main part
- ... which is an isomorphism on the generic fiber,
- and such that $X_1 \rightarrow B_1$ is **semistable**.

Family resolution

- Both theorems answer the question

how well can one resolve a family $X \rightarrow B$ of complex varieties,

with different notions of what you allow to do to B and X , and what you hope to get in the resulting $X' \rightarrow B'$.

- The case $\dim B = 0$ is Hironaka's resolution of singularities.
- The case $\dim B = 1$ is [KKMS]

Definition

A morphism $X_1 \rightarrow B_1$ of smooth complex varieties with $\dim B_1 = 1$ is **semistable** if in local coordinates it is given by

$$t = x_1 \cdots x_k.$$

The one-dimensional base case, and what Carlos said

Theorem (Knudsen-Mumford-Waterman 1973)

Given any family $X \rightarrow B$ with $\dim B = 1$ there is

$$\begin{array}{ccc} X_1 & \rightarrow & X \\ \downarrow & & \downarrow \\ B_1 & \rightarrow & B \end{array}$$

with

- $B_1 \rightarrow B$ an alteration and
- $X_1 \rightarrow X \times_B B_1$ modification,
- such that $X_1 \rightarrow B_1$ is semistable.

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This gives geometric justification for good Hodge theoretic behavior: given $X \rightarrow B$ arbitrary, after base change $B_1 \rightarrow B$ and modification X_1 , the family $X_1 \rightarrow B_1$ has unipotent monodromy at every generic point the discriminant $\Delta(X/B) \subset B$.

What Mumford said

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the Borel-Baily "minimal" compactification). We would also like to study semi-stable reduction over a higher dimensional base: viz., given any dominating morphism $f: X \rightarrow Y$, replacing Y by any Y' generically finite and proper over Y and X by a blow-up of the component of $X \times_Y Y'$ dominating Y' , simplify all the fibres $f': X' \rightarrow Y'$ as much as possible while requiring that X', Y' are non-singular and f is flat.

Toroidal morphisms, log smooth morphisms (KKMS, K. Kato, \aleph -Karu)

- A **toroidal embedding** $U \subset X$ is an open embedding étale locally isomorphic to the embedding of a torus in a toric variety.
- It is the same as a log structure on X **log smooth** over $\text{Spec } k$.
- A **toroidal morphism** between toroidal embeddings $U_i \subset X_i$ is a morphism $X_1 \rightarrow X_2$ that is étale locally the pullback of a dominant toric morphism.
- It is the same as a **log smooth morphism** between log smooth schemes.
- It is characterized by the fact that the pullback of a monomial is a monomial, and is smooth otherwise.
- Once you are log smooth, everything is combinatorial.

Weak toroidalization

The first step is

Theorem (Karu 2000)

Let $X \rightarrow B$ be a dominant morphism of complex varieties. There is a diagram

$$\begin{array}{ccccc} U_X \hookrightarrow & X' & \rightarrow & X & \\ \downarrow & \downarrow & & \downarrow & \\ U_B \hookrightarrow & B' & \rightarrow & B & \end{array}$$

with $B' \rightarrow B$ and $X' \rightarrow B'$ modifications, and $X' \rightarrow B'$ *logarithmically smooth / toroidal*.

- The proof used de Jong's alterations, so could not be made functorial. The generic fiber was modified.

Updated proof, Step 1 (de Jong)

Theorem (Altered semistable reduction, de Jong 1997)

Let $X \rightarrow B$ be a generically smooth complex projective family of varieties. There is a finite group G and a G -equivariant diagram

$$\begin{array}{ccccc} U_Y \hookrightarrow & Y & \rightarrow & X & \\ \downarrow & \downarrow & & \downarrow & \\ U_B \hookrightarrow & B_1 & \rightarrow & B & \end{array}$$

with

- $B_1 \rightarrow B$ and $Y \rightarrow X$ alterations,
- $Y/G \rightarrow X$ and $B_1/G \rightarrow B$ birational,

such that $Y \rightarrow B_1$ is **semistable**.

Consider $\mathcal{X} = [Y/G] \rightarrow X$ and $\mathcal{B} = [B_1/G] \rightarrow B$.

Updated proof, Step 2 (Bergh-Rydh)

Consider $\mathcal{X} = [Y/G] \rightarrow X$ and $\mathcal{B} = [B_1/G] \rightarrow B$. $\mathcal{X} \rightarrow \mathcal{B}$ is log smooth.

Theorem (Destackification, Bergh-Rydh p2019)

There is a diagram

$$\begin{array}{ccccc} X' & \leftarrow & \tilde{\mathcal{X}} & \rightarrow & \mathcal{X} \\ \downarrow & & \downarrow & & \downarrow \\ B' & \leftarrow & \tilde{\mathcal{B}} & \rightarrow & \mathcal{B} \end{array}$$

where $\tilde{\mathcal{X}} \rightarrow \mathcal{X}$ and $\tilde{\mathcal{B}} \rightarrow \mathcal{B}$ are stack blowup sequences, $\tilde{\mathcal{X}} \rightarrow X'$ and $\tilde{\mathcal{B}} \rightarrow B'$ coarse moduli spaces, and $X' \rightarrow B'$ log smooth.

The resulting diagram

$$\begin{array}{ccccc} U_X \hookrightarrow X' & \rightarrow & X & & \\ \downarrow & & \downarrow & & \downarrow \\ U_B \hookrightarrow B' & \rightarrow & B & & \end{array}$$

finishes the proof.

Weakly semistable and semistable morphisms (N-Karu, T. Tsuji)

- A toroidal morphism $X \rightarrow B$ is **weakly semistable** if it is flat with reduced fibers.
- This is the same as an integral and saturated morphism of log structures.
- A toroidal morphism is **semistable** if moreover X and B are smooth.
- In local coordinates, we obtain

$$t_1 = x_1 \cdots x_{l_1}$$

$$\vdots$$

$$t_m = x_{l_{m-1}+1} \cdots x_{l_m}$$

- Yes, this is the best one can get (Karu t1999).

Weak Semistable reduction (Karu)

Theorem (Karu 2000)

Let $X \rightarrow B$ be a generically smooth complex projective family of varieties. There is a diagram

$$\begin{array}{ccccc} U_X & \hookrightarrow & X_1 & \rightarrow & X \\ \downarrow & & \downarrow & & \downarrow \\ U_B & \hookrightarrow & B_1 & \rightarrow & B \end{array}$$

with

- $B_1 \rightarrow B$ an alteration,
- $X_1 \rightarrow (X \times_B B_1)_{\text{main}}$ a modification of the main part
- and such that $X_1 \rightarrow B_1$ is **weakly semistable**.

By weak toroidalization we may assume $X \rightarrow B$ **logarithmically smooth**.

Updated functorial proof (Molcho) part 1

Recall [KKMS] functor:

$$\{\text{toroidal embeddings}\} \xrightarrow{X \mapsto \Sigma_X} \{\text{R.P. cone complexes}\}.$$

It restricts to an equivalence:

$$\{\text{representable tor. modifications}\} \longleftrightarrow \{\text{subdivisions}\}.$$

Theorem (Molcho p2016)

The functor Σ restricts to an equivalence

$$\{\text{stack toroidal modifications}\} \longleftrightarrow \{\text{lattice altered subdivisions}\}.$$

Updated functorial proof (Molcho) part 2

Proposition (Molcho p2016)

$f : X \rightarrow B$ is semistable if and only if $\Sigma(f) : \Sigma_X \rightarrow \Sigma_B$ satisfies

- for all $\sigma \in \Sigma_X$, the image $\Sigma(f)(\sigma)$ is a cone of Σ_B ,
 - for all $\sigma \in \Sigma_X$, the image $\Sigma(f)(N_\sigma) = N_{\Sigma(f)(\sigma)}$.
-
- By the theorem, there is a stack theoretic modification $\mathcal{B} \rightarrow B$ such that the toroidal pullback $\mathcal{X} \rightarrow \mathcal{B}$ is a representable semistable morphism.
 - Using Kawamata's trick one replaces \mathcal{B} by a scheme alteration.

Beyond weak semistable reduction

- In [Karu 2000] we conjectured that **weakly semistable** can be replaced by **semistable**,
- and reduced the problem to polyhedral combinatorics,
- as recently proved by Adiprasito, Liu and Tëmkin.
- This is in parallel to the Knudsen–Mumford–Waterman result.

- We also conjectured that toroidalization can be done more functorially.
- To tell the story we need to go one step back.

Varieties and log structures (K. Kato, Fontaine, Illusie)

- A variety is locally embedded in a smooth variety.
- A log variety is something locally embedded in a toroidal variety.
- The toroidal $U \subset X$ is encoded in the multiplicative submonoid $M_X \subset \mathcal{O}_X$ of functions invertible on U .
- In general a log structure $M \rightarrow \mathcal{O}_Y$ is a morphism of sheaves of monoids inducing an isomorphism on \mathcal{O}_Y^\times .
- A key example is a point on a toric variety.

Resolution and log resolution

- By Hironaka, a variety can be canonically resolved.
- Włodarczyk showed the benefits of **functorial** resolution: if the procedure is functorial for **smooth** morphisms, then gluing and descent is automatic.
- A morphism $X \rightarrow B$ has little chance of having a smooth resolution.
- Toroidalization $[\mathbb{N}K]$ is precisely **log smooth resolution**.
- To make it functorial we turn to Hironaka–Włodarczyk methods.

Functoriality in log resolution

A **logarithmically functorial resolution** assigns to a log morphism $X \rightarrow B$ a modification $X' \rightarrow X$ such that

- $X' \rightarrow B$ is log smooth
- If $Y \rightarrow X$ is log smooth, with log resolution $Y' \rightarrow Y$, then $Y' = Y \times_X^{\log} X' \dots$

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Either that, or it says “sorry, my friend, please modify B first”.

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- A good base change $B' \rightarrow B$ always exists, and
- $X' \rightarrow X$ commutes with base change.

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It is a **relatively log functorial resolution** if moreover

- A good base change $B' \rightarrow B$ always exists, and
- $X' \rightarrow X$ commutes with base change.

The theorem of [NTW 2020] says that a relatively log functorial resolution exists, with the caveat of the next slides.

(We have a draft of a result showing $B' \rightarrow B$ can be made functorial when $X \rightarrow B$ proper.)

Example: log modification of B

- Consider $X = \mathbb{A}_{\log}^1 \rightarrow B = \mathbb{A}^1$.
- It is not log smooth.
- Modify $B' = \mathbb{A}_{\log}^1$,
- then the log pullback $X \rightarrow B'$ is log smooth.

Example: log resolution 1

- Consider $B = \mathbb{A}_{\log}^1$, $Y = \mathbb{A}_{\log}^1 \times \mathbb{A}^1$ and $X = V((x - y)(x + y)) = V(x^2 - y^2)$.
- To resolve X , we blow up the origin (x, y) on Y , **including the exceptional in the log structure.**
- the log proper transform $X' \rightarrow B'$ is log smooth.

Example: log resolution 2

- Consider $B = \mathbb{A}_{\log}^1$, $Y = \mathbb{A}_{\log}^1 \times \mathbb{A}^1$ and $X = V(x_1 - y^2)$.
- Its pullback via $x_1 = x^2$ is example 1.
- By log functoriality we must blow up something whose pullback is (x, y) .
- In other words, we must blow up $(\sqrt{x_1}, y)$.
- This is a **Kummer blow up**, whose result is a stack theoretic blowup.
- the log stack proper transform $\mathcal{X}' \rightarrow B'$ is log smooth.

Example 2 computation

- We blow up $(\sqrt{x_1}, y)$:
- Consider $x_1 = U^2 x'_1$, $y = Uy'$,
- with \mathbb{G}_m action $(x'_1, y', U) \mapsto (t^2 x'_1, ty', t^{-1}U)$.
- The map $\text{Spec } k[x'_1, y', U] \rightarrow \text{Spec } k[x_1, y]$ is \mathbb{G}_m -equivariant,
- leaving $Z := V(x'_1, y')$ invariant.
- Write $\mathcal{X}' = [(\text{Spec } k[x'_1, y', U] \setminus Z)/\mathbb{G}_m]$.

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- leaving $Z := V(x'_1, y')$ invariant.
- Write $\mathcal{X}' = [(\text{Spec } k[x'_1, y', U] \setminus Z)/\mathbb{G}_m]$.
- The equation $x_1 - y^2$ becomes $U^2(x'_1 - y'^2)$,
- and the proper transform $(x'_1 - y'^2)$ is indeed log smooth over \mathbb{A}_{\log}^1 .

Lesson learned

- So log smooth functoriality **requires** log stacks.
- With Bergh's destackification, we get a schematic log resolution as in the theorem,
- which is functorial only for **smooth** $Y \rightarrow X$.
- **more to come tomorrow.**

Also, Hodge theorists,

- One can have, with good reason, monodromy unipotent everywhere,
- with very nice local equations everywhere,
- and functoriality properties.

The end

Thank you for your attention!