

Real Numbers

9/13/07

real # - infinite binary fraction
(decimal)

"Naive" definition

Countably many exceptions (Rational #): $\exists$ more than one representation.

ex) 0.1000...
0.0111...

$$\text{Card}([0,1] \cup \text{exceptions}) = \text{card} \left\{ \{x_n\}_0^\infty : x_n = 0,1 \right\} = \text{card } 2^{\mathbb{N}}$$

collection of all binary sequences

$$2^{\mathbb{N}} \xleftrightarrow{\varphi} \{x_n\}_1^\infty$$

$$\varphi(\{x_n\}) = \{k \in \mathbb{N} : x_k = 1\}$$

prop X - infinite, C countable

$$\Rightarrow \text{Card}(X \cup C) = \text{card}(X)$$

Proof as an exercise

$\leftarrow$ Implies that $\text{Card } \mathbb{R} = \text{card } 2^{\mathbb{N}}$

$$\text{card } 2^{\mathbb{N}} \leq \text{card } 2^{\mathbb{N}} \times 2^{\mathbb{N}} \dots \textcircled{1}$$

$$\begin{array}{ccc} A & \mapsto & A \times A_0 \\ \cap & & \cap \\ \mathbb{N} & & \mathbb{N} \end{array} \quad \begin{array}{l} \text{fixed} \\ \text{(injection)} \end{array}$$

$\leftarrow$ Explanation that
 $\text{Card } \mathbb{R} \times \mathbb{R} =$
 $\text{Card } \mathbb{R}$

Also, claim $2^{\mathbb{N}} \times 2^{\mathbb{N}} \subset 2^{\mathbb{N} \times \mathbb{N}}$

$$\begin{array}{ccc} A \times B & & C \subset \mathbb{N} \times \mathbb{N} \\ \cap & & \cap \\ \mathbb{N} & & \mathbb{N} \end{array}$$

$$\Rightarrow \text{card } 2^{\mathbb{N}} \times 2^{\mathbb{N}} \leq \text{card } 2^{\mathbb{N} \times \mathbb{N}} = \text{card } 2^{\mathbb{N}} \quad \text{because } \mathbb{N} \times \mathbb{N} \text{ countable} \dots \textcircled{2}$$

$$\text{By } \textcircled{1} \text{ \& } \textcircled{2}, \quad \text{card } 2^{\mathbb{N}} \times 2^{\mathbb{N}} = \text{card } 2^{\mathbb{N} \times \mathbb{N}} = 2^{\mathbb{N}}$$

$\mathbb{R}$ - complete ordered field

Formal def of $\mathbb{R}$

(
↳ operations $+, -, \times, \div$ (satisfying properties see textbook)
↳ total order (only three possibilities: $\alpha > \beta$, $\beta > \alpha$ or $\alpha = \beta$)
& agree with algebraic operations.

$$\text{if } a \geq b, c \geq d \Rightarrow a+c \geq b+d$$

$$a > b, c > 0 \Rightarrow ac > bc$$

• Finite field cannot be ordered

• $\mathbb{C}$ cannot be ordered.

Complete

Let set $A \subset \mathbb{R}$ (ordered field).

A is bounded above, if $\exists M$ s.t. $\forall x \in A, x \leq M$

Def) $\sup A$.

The smallest upper bound of A (if exists) is called supremum of A .

$\sup A = M$ s.t. M is upperbound for A , ^{but} $\forall m < M, m \in \mathbb{R}$ is not upper bound.

* Complete means any ^{non-empty} bdd above A has supremum.

ex) ①. $A = \{x \in \mathbb{Q} : x^2 < 2\}$ A does not have a rational supremum.

$\sup A = \sqrt{2} \Rightarrow \mathbb{Q}$ is not complete.

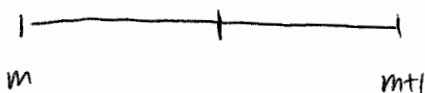
Why $\mathbb{R}$ is complete. (If $\mathbb{R}$ is set of binary fractions)

$\mathbb{R}$ - binary fractions (infinite)

$A \subset \mathbb{R}, A \cap (0, \infty) \neq \emptyset \leftarrow$ Assume for simplicity

Let $m = \min \{n : \forall x \in A, x < n+1\}$ m gives us integer part of $\sup$.

$m, b_1 b_2 b_3 \dots$



We can find inductively all binary digits b_i of $\sup A$:

If $A \cap [m + \frac{1}{2}, m+1) = \emptyset$ then $b_1 = 0$
 $\neq \emptyset$ then $b_1 = 1$

$A \cap [m + \frac{1}{2} b_1, m + \frac{1}{2} b_1 + \frac{1}{4}) = \emptyset$ then $b_2 = 0$
 $\neq \emptyset$ then $b_2 = 1$
 $\vdots$

$A \cap [m + \frac{1}{2} b_1 + \frac{1}{4}, m + \frac{1}{2} b_1 + \frac{1}{4} + \frac{1}{8}) = \emptyset$
 next step: check
 $A \cap [m + \frac{1}{2} b_1 + \frac{1}{4} b_2, m + \frac{1}{2} b_1 + \frac{1}{4} b_2 + \frac{1}{8}) = \emptyset$

Thm There exists a "unique" complete ordered field.

If $\exists \mathbb{F}_1, \mathbb{F}_2$ both complete ordered field. $\leftarrow$ "Unique" means
 $\Rightarrow \exists \varphi: \mathbb{F}_1 \rightarrow \mathbb{F}_2$ preserving algebraic ops & order.
 bijection

Dedekind cuts (Earliest formal construction of $\mathbb{R}$)

Cut: subset $I \subset \mathbb{Q}$ s.t. 1. I is bounded above

2. If $x \in I, y < x \Rightarrow y \in I$

3. I has no max element.

4. $I \neq \emptyset$
 $I \neq \mathbb{Q}$

$\mathbb{Q} \cap (-\infty, x) \leftarrow$ Informally, Ded. cut is

If we know cut, we know all $q \in \mathbb{Q}, q < x$. So we know all $r \in \mathbb{Q}, r > x$

Therefore, we can approximate x by rationals with any given precision

Completeness of $\mathbb{R}$ in this def is easy.

If we have a collection A of cuts

then $\bigcup_{\tilde{x} \in A} \tilde{x} \quad (\subset \mathbb{Q})$

is the cut corresponding to $\sup$

(Each $\tilde{x}$ is a subset of $\mathbb{Q}$, so is the union)

See s. 1.9 for details