

Ergodicity of Markov processes: theory and computation

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- 1 Markov processes on measurable state space.
- 2 Coupling method and renewal theory
- 3 Exponential and power-law ergodicity
- 4 Construction of Lyapunov functions
- 5 Numerical computation of ergodicity
- 6 Numerical computation of invariant probability measures

Basic setting 1

- 1 Φ_n – discrete time Markov process
- 2 $(X, \mathcal{B}(X))$ – state space with a sigma algebra $\mathcal{B}(X)$
- 3 P – transition probability. $P(x, A) = \mathbb{P}[\Phi_1 \in A \mid \Phi_0 = x]$.
- 4 $P(x, \cdot)$ is a probability measure on $(X, \mathcal{B}(X))$, $P(x, A)$ is a measurable function for any $A \in \mathcal{B}$.
- 5 By Markov property, this is enough to determine a Markov process

Markov property: only depends on the nearest history

$$\mathbb{P}[\Phi_{n+1} \in A \mid \Phi_0, \dots, \Phi_n] = \mathbb{P}[\Phi_{n+1} \in A \mid \Phi_n]$$

- $P^m(x, A) = \mathbb{P}[\Phi_{n+m} \in A \mid \Phi_n = x]$.



$$P^{m+n}(x, A) = \int_{\mathcal{X}} P^n(y, A) P^m(x, dy)$$

- First arrival time: $\eta_A = \inf_{n \geq 1} \{\Phi_n \in A\}$
- Note that η_A is a stopping time (random time that only depends on historical and present states of Φ_n .)
- Hitting probability: $L(x, A) = \mathbb{P}[\Phi_n \in A \text{ for some } n \mid \Phi_0 = x]$

Main difference from discrete Markov chain: $P(x, y)$ does not make sense any more!

Φ_n is irreducible if there exists a reference measure ψ on X such that

- 1 If $\psi(A) > 0$, then $L(x, A) > 0$ for all $x \in X$
- 2 If $\psi(A) = 0$, then $\psi(\{y : L(y, A) > 0\}) = 0$

Φ_n can reach everywhere that could be “seen” by ψ .

Example

Stochastic differential equation X_t . Euler-Maruyama method.

$$X_{n+1} = X_n + f(X_n)h + \sigma(X_n)\mathcal{N}(0, 1)\sqrt{h}$$

Transition kernel

$$P(x, A) = \int_A \frac{1}{\sqrt{2\pi\sigma(x)^2h}} e^{-(y-x-f(x)h)^2/2\sigma^2(x)h} dy$$

Let Lebesgue measure be the reference measure. Easy to check that X_n is irreducible.

Atom and pseudo-atom

- 1 Discrete state space: $P(x, y) > 0$. Very useful!
- 2 Atom: α is an atom if $P(x, \cdot) = P(y, \cdot)$ for all $x, y \in \alpha$. Atom is like a discrete state.
- 3 Atom usually does not exist
- 4 Pseudo-atom: small set C
- 5 $C \in \mathcal{B}(X)$ is a small set if there exist an integer $n \in \mathbb{N}$ and a nontrivial measure ν such that

$$P^n(x, A) \geq \nu(A) \text{ for all } x \in C$$

Example

Euler-Maruyama scheme again

$$X_{n+1} = X_n + f(X_n)h + \sigma(X_n)\mathcal{N}(0, 1)\sqrt{h}$$

Every bounded set is a small set because the probability density of P is everywhere strictly positive.

Random walk: $X_{n+1} = X_n + U_n$, $U_n \sim U(-1/2, 1/2)$.
 $[-1/4, 1/4]$ is a small set with $n = 1$ and $\nu = \text{Lebesgue measure}$.

(A)periodicity

Discrete space

Assume irreducibility. Define $E = \{n \mid P^n(x, x) > 0\}$. Period d is the greatest common divisor of E .

General space

Assume irreducibility. C is a small set. Define

$$E_C = \{n \mid P^n(x, \cdot) \geq \nu(\cdot), x \in C, \nu(C) > 0\}$$

(positive probability that the chain will return to C after n steps.)
Period d is the greatest common divisor of E .

Φ_n is aperiodic if $d = 1$.

From now on we assume that Φ_n is irreducible and aperiodic.

- 1 Left operator: μ -probability measure. $\mu P^n(A) = \mathbb{P}_\mu[\Phi_n \in A]$.
- 2 Right operator: f -observable (function). $P^n f(x) = \mathbb{E}_x[f(\Phi_n)]$.
- 3 Invariant probability measure. π is said to be invariant if $\pi P = \pi$.

Let μ and ν be two probability measures. Does

$$\|\mu P^n - \nu P^n\|_{TV}$$

converge to zero? If yes, how fast??

Main approach: Coupling

A Markov process (Φ_n^1, Φ_n^2) on the state space $X \times X$ is said to be a Markov coupling if

- 1 Two marginal distributions are Markov processes Φ_n with initial distribution μ and ν , respectively
- 2 If $\Phi_n^1 = \Phi_n^2$, then $\Phi_m^1 = \Phi_m^2$ for all $m \geq n$.

$\tau_C = \inf_{n \geq 0} \{\Phi_n^1 = \Phi_n^2\}$ is the *coupling time*.

Coupling Lemma

$$\|\mu P^n - \nu P^n\|_{TV} \leq 2\mathbb{P}[\tau_C > n].$$

(See whiteboard for the proof.)

Optimal coupling (Pitman 1970s)

There exists a coupling (Φ_n^1, Φ_n^2) (may not be Markov) such that

$$\|\mu P^n - \nu P^n\|_{TV} = 2\mathbb{P}[\tau_C > n].$$

The existence of “honest” optimal coupling remains open.

Coupling at atom

- 1 Assume Φ_n admits an atom α .
- 2 Let (Φ_n^1, Φ_n^2) be a coupling such that Φ_n^1 and Φ_n^2 are independent until their first simultaneous visit to α , and run together after that.

Easy to check: (Φ_n^1, Φ_n^2) is a Markov coupling.

Difficulty: property of $\mathbb{P}[\tau_C > n]$?

- 1 Exponential: $\mathbb{P}[\tau_C > n] \sim \rho^{-n}$ for $\rho > 1$
- 2 Power-law: $\mathbb{P}[\tau_C > n] \sim n^{-\beta}$ for $\beta > 0$

Let

$$S_n = \sum_{i=0}^n Y_i$$

such that Y_1, Y_2, \dots are i.i.d. random nonnegative integers. (Y_0 could be different). S_n is a renewal process. Y_i is called inter-occurrence time.

Let $u_n = \mathbb{P}[n = S_m \text{ for some } m]$.

If S is aperiodic, $u_n \rightarrow 1/\mathbb{E}[Y_1]$.

Renewal process from Φ_n

- 1 α is the atom.
- 2 $Y_0 = \eta_\alpha$
- 3 S_n is the n -th visit to α
- 4 S_n is a renewal process because α is an atom. $Y_i = \eta_\alpha | \Phi_0 = \alpha$.
(Markov property: history is independent of the future.)

Simultaneous renewal

- 1 Now let S_n and S'_n be two renewal processes corresponding to Φ_n^1 and Φ_n^2 , respectively.
- 2 The coupling time τ_C is the first simultaneous renewal time.

$$\tau_C = \inf_n \{n = S_{k_1} = S'_{k_2} \text{ for some } k_1 \text{ and } k_2\}$$

Three questions

- 1 What if there is no atom? ✓
- 2 First simultaneous renewal time? ✓
- 3 How to estimate the first visit time η_α (probably tomorrow)

How to make an atom? (1)

- 1 Atom does not exist in most scenarios
- 2 Small set is much easier to get
- 3 Simplest case. Let C be a small set that satisfies

$$P(x, A) \geq \delta \mathbf{1}_C(x) \nu(A) \quad , \quad A \in \mathcal{B}(X), x \in X,$$

where ν is a probability measure with $\nu(C) = 1$.

- 4 Split X into $\hat{X} = X \times \{0, 1\}$ with $X_0 = X \times \{0\}$ and $X_1 = X \times \{1\}$.
- 5 Similarly, split A into A_0 and A_1

How to make an atom? (2)

- ① Let λ be a measure on X . Split λ into $\hat{\lambda}$ on \hat{X} such that

$$\lambda^*(A_0) = \lambda(A \cap C)(1 - \delta) + \lambda(A \cap C^c)$$

$$\lambda^*(A_1) = \lambda(A \cap C)\delta$$

- ② In other words, $\lambda^*(A_0 \cup A_1) = \lambda(A)$
- ③ Split transition kernel P into \hat{P} :

$$\hat{P}(x, \cdot) = P(x, \cdot)^* \quad x \in X_0 \setminus C_0$$

$$\hat{P}(x, \cdot) = (1 - \delta)^{-1}[P(x, \cdot)^* - \delta\nu^*(\cdot)] \quad x \in C_0$$

$$\hat{P}(x, \cdot) = \nu^*(\cdot) \quad x \in C_1$$

How to make an atom? (3)

- 1 A Markov process $\hat{\Phi}_n$ is defined on \hat{X} with transition probability \hat{P} .
- 2 C_1 becomes an atom.
- 3 Most result (irreducibility, aperiodicity, recurrence etc.) still holds for $\hat{\Phi}_n$

First simultaneous renewal time?

- 1 $S_n = Y_0 + Y_1 + \cdots + Y_n$, $S'_n = Y'_0 + Y'_1 + \cdots + Y'_n$
- 2 $Y_0 = \eta_\alpha | \Phi_0 \sim \mu$, $Y'_0 = \eta_\alpha | \Phi_0 \sim \nu$
- 3 $Y_1, Y'_1, Y_2, Y'_2, \dots$ are i.i.d. with distribution $\eta_\alpha | \Phi_0 = \alpha$
- 4 Let T be the simultaneous renewal time

$$T = \inf_n \{n = S_{k_1} = S'_{k_2} \text{ for some } k_1, k_2\}$$

- 5 From renewal theorem: There exist n_0 and c such that

$$\mathbb{P}[n \text{ is a renewal time}] = \mathbb{P}[n = S_k \text{ for some } k] \geq c$$

for all $n \geq n_0$.

Exponential tail

If $\mathbb{E}[\rho_1^{Y_0}], \mathbb{E}[\rho_1^{Y'_0}], \mathbb{E}[\rho_1^{Y_1}] < \infty$ for some $\rho_1 > 1$, then there exists $\rho_0 > 1$ such that $\mathbb{E}[\rho_0^T] < \infty$.

Power-law tail

If $\mathbb{E}[Y_0^\beta], \mathbb{E}[(Y'_0)^\beta], \mathbb{E}[Y_1^\beta] < \infty$ for some $\beta > 0$, then $\mathbb{E}[T^\beta] < \infty$.

(Note that finite exponential/power-law moment is equivalent to exponential/power-law tail.)

Proof on whiteboard.

Ref: Lectures on the Coupling Method by Torgny Lindvall

Thank you